

# Kullback - Leibler divergence

Short talk - 04/10/2021

# KL expression



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- Discrete version:

$$D_{KL}(p\|q) = \sum_{x \in \mathcal{X}} p(x) \log \frac{p(x)}{q(x)}$$

- Continuous version:

$$D_{KL}(p\|q) = \int_{-\infty}^{+\infty} p(x) \log \frac{p(x)}{q(x)} dx$$

If there exists  $x \in \mathcal{X}$  such that  $q(x) = 0$  and  $p(x) \neq 0$  then  $D(p\|q) = +\infty$

$$\log \frac{p(x)}{q(x)}$$

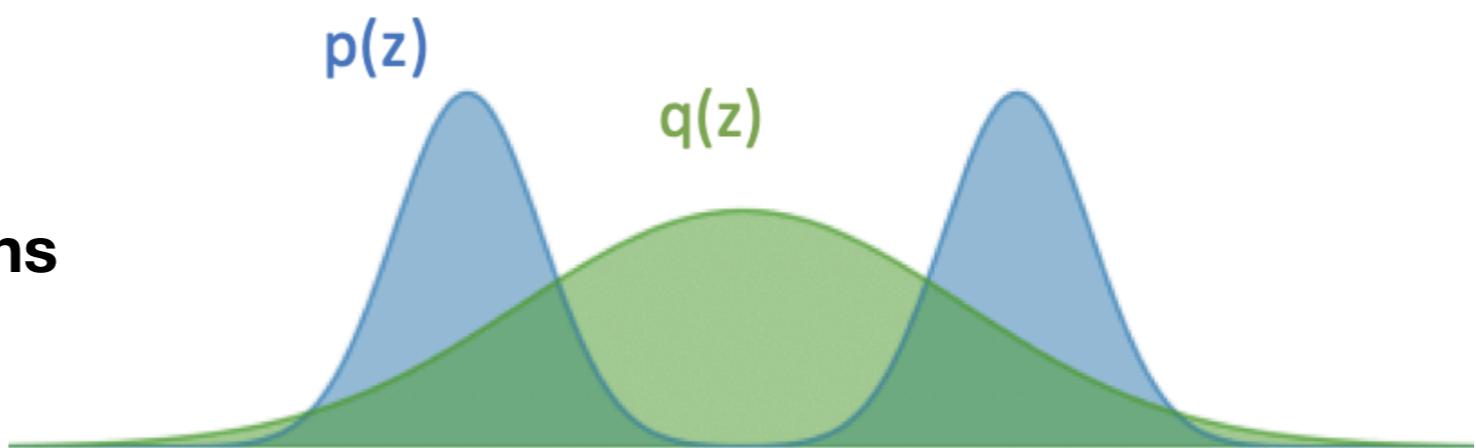
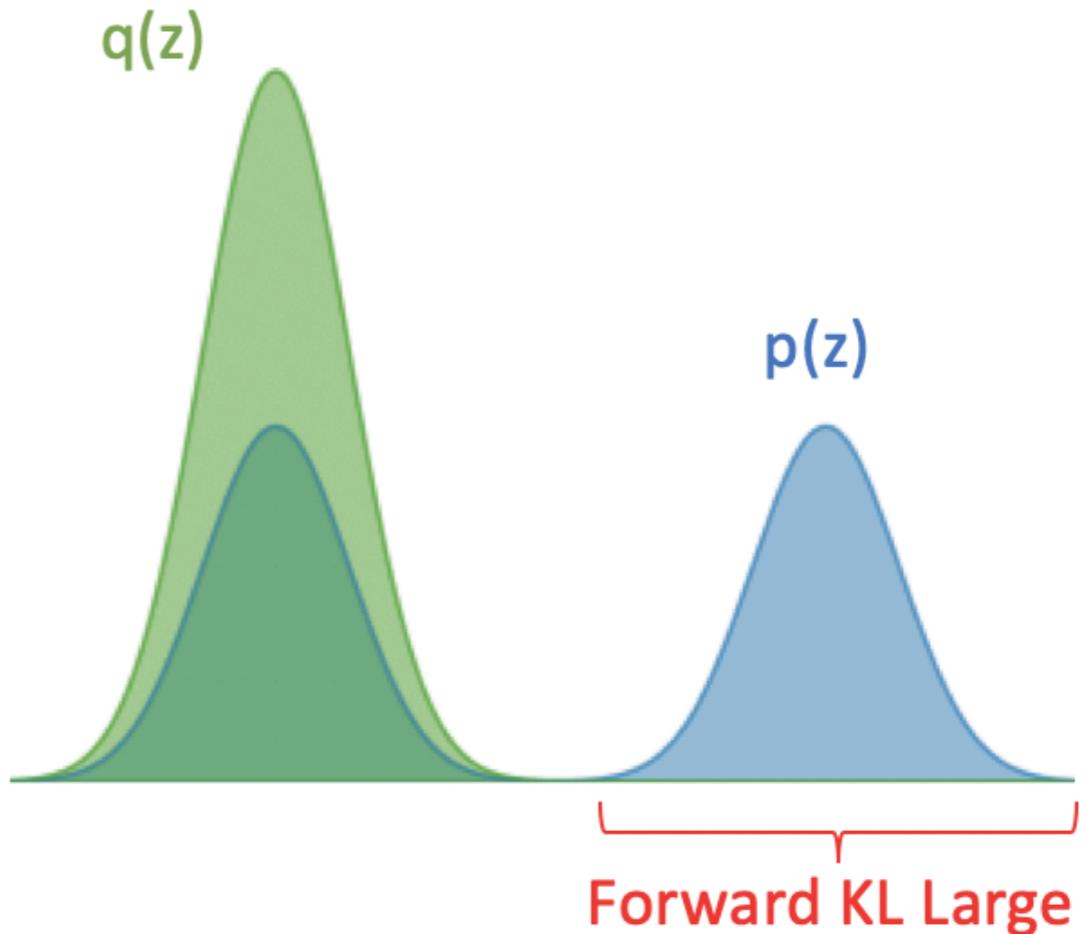
- Positive if  $p(x) \geq q(x)$
- Null if  $p(x) = q(x)$
- Negative if  $q(x) \geq p(x)$

**Penalties are weighted by  $p$  :**

$$D_{KL}(p\|q) = \sum_{x \in \mathcal{X}} p(x) \log \frac{p(x)}{q(x)}$$

**Keep in mind that the above functions  
are probability densities:**

$$\sum_{x \in \mathcal{X}} p(x) = 1 \quad \text{and} \quad \sum_{x \in \mathcal{X}} q(x) = 1$$



**$D_{KL}(p\|q) \geq 0$  and equality holds if  $p = q$**

$$\begin{aligned} -D_{KL}(p\|q) &= - \sum_{x \in \mathcal{X}} p(x) \log \frac{p(x)}{q(x)} = \sum_{x \in \mathcal{X}} p(x) \log \frac{q(x)}{p(x)} \\ &\leq \log \sum_{x \in \mathcal{X}} p(x) \frac{q(x)}{p(x)} = 0 \quad (\textbf{Jensen inequality}) \end{aligned}$$

**$D_{KL}(p\|q) = 0$  if equality in Jensen inequality i.e.**

**$p = cq \rightarrow p = q$  since  $p$  and  $q$  sum to 1.**

# Definition of a distance

Any function  $d : \mathcal{X} \times \mathcal{X} \rightarrow \mathbb{R}_+$  such that :

- $\forall (a, b) \in \mathcal{X}^2, \quad d(a, b) = 0 \iff a = b$



- $\forall (a, b) \in \mathcal{X}^2, \quad d(a, b) = d(b, a)$



- $\forall (a, b, c) \in \mathcal{X}^3, \quad d(a, c) \leq d(a, b) + d(b, c)$



**$D_{KL}(p||q)$  is not a distance !**

## Information theory intuition :

$$D_{KL}(p\|q) = \underbrace{\sum_{x \in \mathcal{X}} p(x) \log p(x)}_{H(p)} - \underbrace{\sum_{x \in \mathcal{X}} p(x) \log q(x)}_{H_p(q)}$$

If we knew the true distribution  $p$  of the random variable, we could construct a code with average description length  $H(p)$ .

If, instead, we used the code for a distribution  $q$ , we would need  $H(p) + D_{KL}(p\|q)$  bits on the average to describe the random variable.

**Let  $x_1, \dots, x_N \in \mathcal{X}$  be  $N$  i.i.d. observations of a random variable  $X$**

$$\hat{p}(x) = \frac{1}{N} \sum_{i=1}^N \delta(x - x_i)$$

**Let  $p_\theta$  be a parameterized distribution on  $\mathcal{X}$**

$$D_{KL}(\hat{p} \| p_\theta) = \sum_{x \in \mathcal{X}} \hat{p}(x) \log \frac{\hat{p}(x)}{p_\theta(x)} = -H(\hat{p}) - \sum_{x \in \mathcal{X}} \hat{p}(x) \log p_\theta(x)$$

$$= -H(\hat{p}) - \sum_{x \in \mathcal{X}} \sum_{i=1}^N \delta(x - x_i) \log p_\theta(x) = -H(\hat{p}) - \frac{1}{N} \sum_{i=1}^N \log p_\theta(x_i)$$

**Maximizing the likelihood  $p_\theta(x)$   $\iff$  Minimizing  $D_{KL}(\hat{p} \| p_\theta)$**

# Pinsker's inequality

Total variation distance

$$\delta(p, q) = \sup_{A \in \mathcal{F}} |p(A) - q(A)|$$

$$\delta(p, q) \leq \sqrt{\frac{1}{2} D_{KL}(p \| q)}$$

# In practice

**Let  $x_i$  be samples from  $p(x)$  :**

$$\lim_{N \rightarrow +\infty} \frac{1}{N} \sum_{i=0}^N \log \frac{p(x_i)}{q(x_i)} = \sum_{x \in \mathcal{X}} p(x) \log \frac{p(x)}{q(x)} = D_{KL}(p || q)$$